Women in Logic

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MONOIDAL WIDTH

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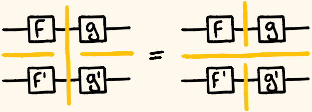
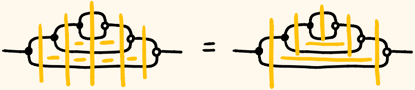
MOTIVATION (1)

- how efficient is to compute the semantics of morphisms in monoidal categories?





- we need an 'algebra of decompositions'

MOTIVATION (2)

- existing notions of complexity for graphs are based on decompositions: path width, tree width, branch width and rank width

- make explicit the algebra of decomposition that is hidden behind the definitions of these graph widths

MAIN RESULTS

- monoidal width as a measure of complexity for morphisms in monoidal categories
- monoidal decomposition as explicit decomposition algebra
- capture some known measures of complexity for graphs: path width, tree width, branch width and rank width

OUTLINE

- monoidal decompositions
- monoidal width for matrices

DECOMPOSITION SYSTEM

A decomposition system $(\mathcal{A}, \mathcal{O}, w)$ in a monoidal category $\mathcal{C}$ is given by

- $\mathcal{A}$: set of 'atomic' morphisms in $\mathcal{C}$
- $\mathcal{O} = \{\otimes, ;_X \text{ for } X \in obj(\mathcal{C})\}$: set of operations
- $w : \mathcal{A} \cup \mathcal{O} \to \mathbb{N}$: weight function
  such that:

$$\begin{cases} w(\otimes) = 0 \\ w(;_{X \otimes Y}) = w(;_X) + w(;_Y) \end{cases}$$

DECOMPOSITION SYSTEM - EXAMPLE

a decomposition system $(\mathcal{A}, \mathcal{O}, w)$

in $\mathcal{C}$ $\rightsquigarrow$ $\mathcal{F}in\mathcal{S}et$

- $\mathcal{A}$: set of 'atoms' $\rightsquigarrow$ $\{\exists\!-, \bullet\!-, \times, -\}$

- $\mathcal{O} = \{\otimes, ;_X \text{ for } X \in obj(\mathcal{C})\}$: set of operations

- $w : \mathcal{A} \cup \mathcal{O} \longrightarrow \mathbb{N}$: weight such that:

$$\begin{cases} w(\otimes) = 0 \\ w(;_{X \otimes Y}) = w(;_X) + w(;_Y) \end{cases}$$

$\rightsquigarrow$ $w(\exists\!-) = w(\times) = 2$
$w(\bullet\!-) = w(-) = 1$
$w(;_n) = n$

MONOIDAL DECOMPOSITION

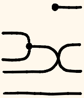
$f : X \to Y$ morphism in $\mathcal{C}$

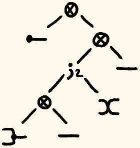
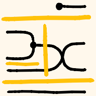
a monoidal decomposition $d \in \mathcal{D}_f$ of f is

$$
\begin{aligned}
d ::= \;& (f) && \text{if } f \in \mathcal{A} \\
\mid \;& d_1 \,{}^{\swarrow\, ;_c \,\searrow}\, d_2 && \text{if } f = f_1 \,;_c\, f_2 \;,\; d_1 \in \mathcal{D}_{f_1} ,\; d_2 \in \mathcal{D}_{f_2} \\
\mid \;& d_1 \,{}^{\swarrow\, \otimes \,\searrow}\, d_2 && \text{if } f = f_1 \otimes f_2 \;,\; d_1 \in \mathcal{D}_{f_1} ,\; d_2 \in \mathcal{D}_{f_2}
\end{aligned}
$$

$\rightsquigarrow$ it's a labelled binary tree

MONOIDAL DECOMPOSITION - EXAMPLE

 $: 4 \longrightarrow 4$ morphism in FinSet

$d =$  $\rightsquigarrow$ 

MONOIDAL WIDTH

$d \in \mathcal{D}_f$ monoidal decomposition of f

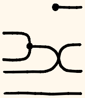
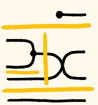
WIDTH OF d

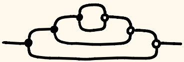
$$
\begin{aligned}
\mathrm{wd}(d) := \;& w(f) && \text{if } d = (f) \\
\mid \;& \max\{\mathrm{wd}(d_1),\, w(;_c),\, \mathrm{wd}(d_2)\} && \text{if } d = \begin{array}{ccc} & ;_c & \\ d_1 \diagup & & \diagdown d_2 \end{array} \\
\mid \;& \max\{\mathrm{wd}(d_1),\, \mathrm{wd}(d_2)\} && \text{if } d = \begin{array}{ccc} & \otimes & \\ d_1 \diagup & & \diagdown d_2 \end{array}
\end{aligned}
$$

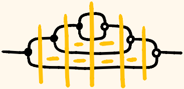
MONOIDAL WIDTH OF f

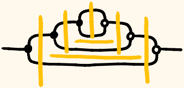
$$\mathrm{mwd}(f) := \min_{d \in \mathcal{D}_f} \mathrm{wd}(d)$$

MONOIDAL WIDTH - EXAMPLE

$f =$  $\mathrm{wd}\Big($  $\Big) = 2$

$4 =$ 

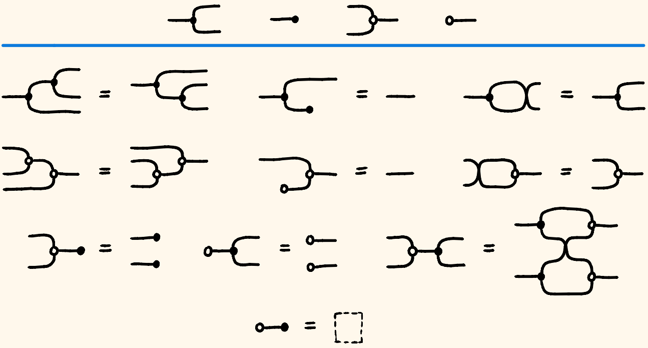
$\mathrm{wd}\Big($  $\Big) = 4$

$\mathrm{wd}\Big($  $\Big) = 2$

OUTLINE

- monoidal decompositions
- monoidal width for matrices

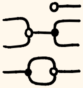
PROP OF MATRICES



[1] Zanasi, Interacting Hopf algebras, PhD thesis (2018)

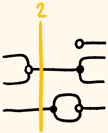
PROP OF MATRICES - EXAMPLE

$$A = \begin{pmatrix} 0 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 2 \end{pmatrix} =$$

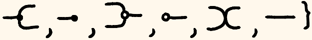


FACT : the minimal vertical cut in a matrix is its rank : $\min\{k \in \mathbb{N} \mid A = B ;_k C\} = \mathrm{rank}\, A$

$\mathrm{rank}\, A = 2 \rightsquigarrow$



MONOIDAL WIDTH OF MATRICES

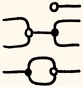
$\mathcal{A} = \{$  $\}$

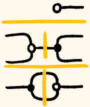
$$A = \begin{pmatrix} A_1 & 0 & \cdots & 0 \\ 0 & A_2 & & \vdots \\ \vdots & & \ddots & \\ 0 & \cdots & & A_b \end{pmatrix} = A_1 \oplus A_2 \oplus \cdots \oplus A_b$$

THEOREM

$$\max_i \operatorname{rank} A_i \leq \operatorname{mwd} A \leq \max_i \operatorname{rank} A_i + 1$$

MONOIDAL WIDTH OF MATRICES - EXAMPLE

$$A = \begin{pmatrix} 0 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 2 \end{pmatrix} =$$ 

$\mathrm{wd}\Big($  $\Big) = 2$

$$= \max\{\underbrace{\mathrm{rank}(\ )}_{=0},\ \underbrace{\mathrm{rank}\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}}_{=1},\ \underbrace{\mathrm{rank}(2)}_{=1}\} + 1$$

SUMMARY OF RESULTS

MATRICES
$$\max_i \operatorname{rank} A_i \leq \operatorname{mwd} A \leq \max_i \operatorname{rank} A_i + 1$$

COSPANS OF GRAPHS
$$\operatorname{pwd}(G) = \operatorname{mpwd}(g)$$
$$\operatorname{twd}(G) \leq \operatorname{mtwd}(g) \leq 2 \cdot \operatorname{twd}(G)$$
$$\tfrac{1}{2} \operatorname{bwd}(G) \leq \operatorname{mwd}(g) \leq \operatorname{bwd}(G) + 1$$

PROP OF GRAPHS
$$\tfrac{1}{2} \operatorname{rwd}(G) \leq \operatorname{mwd}(g) \leq 2 \operatorname{rwd}(G)$$

FUTURE WORK

- obtain a result similar to Courcelle's theorem
- capture other widths (clique width, twin width, ...
  tree width for directed graphs and relational structures)
- algorithmic applications

SOME REFERENCES

[1] Zanasi, Interacting Hopf algebras, PhD thesis (2018)

[2] Oum & Seymour, Approximating clique width and branch width (2006)

[3] Chantawibul & Sobociński, Towards compositional graph theory (2015)

- Bonchi, Piedeleu, Sobociński & Zanasi, Graphical affine algebra (2019)
- Di Lavore, Hedges & Sobociński, Compositional modelling of network games (2021)

THIS WORK

- Di Lavore & Sobociński, Monoidal width: capturing rank width (2022)
- Di Lavore & Sobociński, Monoidal width: unifying tree width, path width and branch width (2022)

RANK WIDTH [Oum & Seymour, 2006]

$G = (V, E, \mathrm{ends} : E \to \mathcal{P}_{\leq 2}(V))$ undirected graph

RANK DECOMPOSITION

(Y, π) where

- Y is a subcubic tree (= any node has at most 3 neighbours)
- $\pi : \mathrm{leaves}\, Y \xrightarrow{\simeq} V$ labelling bijection

WIDTH OF (Y, π)

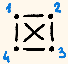
$$wd(Y, \pi) := \max_{e \in \mathrm{edges}\, Y} \mathrm{rank}(X_e)$$

X_e adjacency matrix of the cut given by e through π

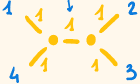
RANK WIDTH

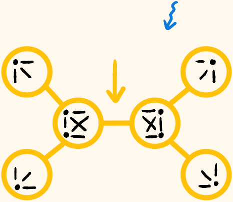
$$rwd(G) := \min_{(Y, \pi)} wd(Y, \pi)$$

RANK WIDTH - EXAMPLE

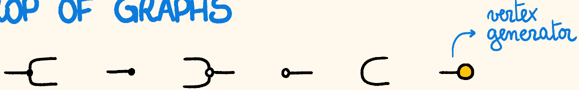
$G =$ 

$$\begin{array}{c|cc} & 2 & 3 \\ \hline 1 & 1 & 1 \\ 4 & 1 & 1 \end{array} \qquad \begin{array}{c|ccc} & 1 & 3 & 4 \\ \hline 2 & 1 & 1 & 1 \end{array}$$

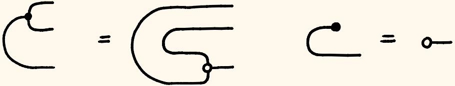




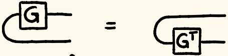
A PROP OF GRAPHS



bialgebra +



⇝ the cup transposes



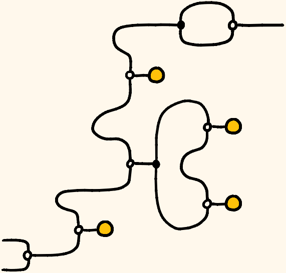
and captures equivalence of adjacency matrices

$G + G^T = H + H^T \iff$

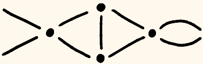
[3] Chantawibul & Sobociński, Towards compositional graph theory (2015)

A PROP OF GRAPHS – EXAMPLE

$\rightsquigarrow$ graph on k vertices given by the adjacency matrix $[G]$



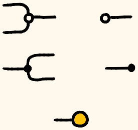
$\rightsquigarrow$



DECOMPOSITIONS IN THE PROP OF GRAPHS

Bialgebra structure

+ 'vertex' generator



ATOMS

$\mathcal{A} = \{\text{all morphisms}\}$

WEIGHT FUNCTION

$w(g) := |\text{vertices } g|$

$w(j_n) := n$

RANK WIDTH & MONOIDAL WIDTH

[G] undirected graph

$g =$  $: 0 \to 0$ in Cgraph

THEOREM

$$\frac{1}{2}\,\mathrm{rwd}(G) \le \mathrm{mwd}(g) \le 2\,\mathrm{rwd}(G)$$