

SYCO 8 - Tallinn

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# MONOIDAL WIDTH

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# MOTIVATION

- measure of complexity for morphisms in a monoidal category
- find a common general framework for different measures of complexity for graphs

# MAIN IDEA

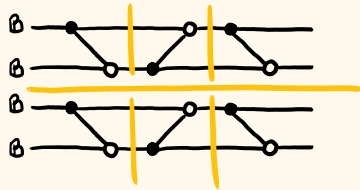
- bring the ideas from graph complexity to monoidal categories
- measure the cost of decomposing a morphism into basic building blocks using some chosen operations

# MAIN IDEA : EXAMPLE

generators =  { CNOT gate }

Operations = { composition, monoidal product }

A monoidal decomposition of exchanging the values of two pairs of variables.



# MAIN RESULTS

- definition of monoidal width for measuring complexity of morphisms
- recover some known notions of complexity for graphs: path width, tree width, branch width, rank width

# OUTLINE

- monoidal decompositions
- monoidal width for matrices
- monoidal width for graphs

# DECOMPOSITION SYSTEMS

$(\mathcal{C}, \mathcal{G}, \mathcal{O}, w)$  decomposition system if

- $\mathcal{C}$  monoidal category
- $\mathcal{G}$  set of generators of  $\mathcal{C}$
- $\mathcal{O} = \{\otimes, ;_x \text{ for any object } x\}$  set of operations
- $w : \mathcal{G} \cup \mathcal{O} \rightarrow \mathbb{N}$  weight function  
such that

$$\begin{cases} w(\otimes) = 0 \\ w(;_{x \otimes y}) = w(;_x) + w(;_y) \end{cases}$$

# DECOMPOSITION SYSTEMS - EXAMPLE

$(\mathcal{C}, \mathcal{F}, \theta, w)$

- $\mathcal{C}$  monoidal category  $\rightsquigarrow \mathcal{C} := \text{FinSet}$
- $\mathcal{F}$  set of generators of  $\mathcal{C}$   $\rightsquigarrow \mathcal{F} := \{-, \int, \circ, \chi\}$
- $\theta = \{\otimes, ;_x \text{ for any object } x\}$  set of operations
- $w : \mathcal{F} \cup \theta \rightarrow \mathbb{N}$  weight  $\rightsquigarrow$   
such that

$$\begin{cases} w(\otimes) = 0 \\ w(;_{x \otimes y}) = w(;_x) + w(;_y) \end{cases}$$

$$\begin{aligned} w(-) &:= 1 \\ w(\int) &:= 2 \\ w(\circ) &:= 1 \\ w(\chi) &:= 2 \\ w(;_m) &:= m \end{aligned}$$



# MONOIDAL DECOMPOSITIONS

$f: X \rightarrow Y$  morphism in  $\mathcal{C}$

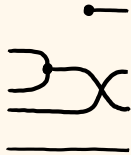
$(S, \mu)$  monoidal decomposition of  $f$  if

- $S$  is a binary tree

- $\mu: \begin{cases} \text{internal nodes}(S) \rightarrow \mathcal{O} \\ \text{leaves}(S) \rightarrow \mathcal{F} \end{cases}$  labelling of  $S$

- $f$  is obtained by assembling the generators on the leaves of  $S$  according to the operations on the internal nodes of  $S$

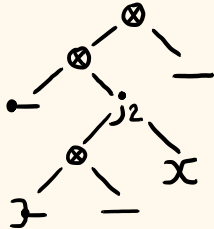
# MONOIDAL DECOMPOSITIONS - EXAMPLE



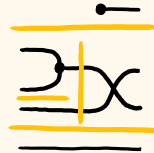
$$: 4 \rightarrow 4$$

morphism in  $\text{FinSet}$

$$(\mathcal{S}, \mu) =$$



$\rightsquigarrow$



# MONOIDAL WIDTH

$(S, \mu)$  monoidal decomposition  
of  $f: X \rightarrow Y$

WIDTH OF A MONOIDAL DECOMPOSITION

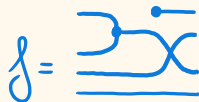
$$\text{wd}(S, \mu) := \max_{m \in \text{modes}(S)} w(\mu(m))$$

MONOIDAL WIDTH

$$\text{mwd}(f) := \min_{(S, \mu)} \text{wd}(S, \mu)$$

# MONOIDAL WIDTH - EXAMPLE

$(S, \mu)$  monoidal decomposition  
of  $f: X \rightarrow Y$

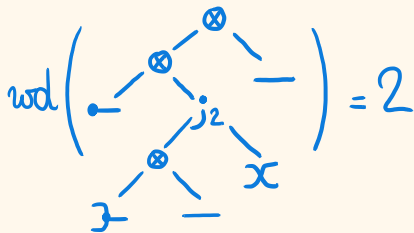


WIDTH OF A MONOIDAL DECOMPOSITION

$$\text{wd}(S, \mu) := \max_{m \in \text{modes}(S)} w(\mu(m))$$

MONOIDAL WIDTH

$$\text{mwd}(f) := \min_{(S, \mu)} \text{wd}(S, \mu)$$



# OUTLINE

- monoidal decompositions

- monoidal width for matrices

- monoidal width for graphs

# MONOIDAL WIDTH OF MATRICES

$$\mathfrak{M} = \{ -, \times, \supset, \circ, \lrcorner, \rightarrow \}$$

$$A = \begin{pmatrix} A_1 \circledast & \dots & \circledast \\ \circledast & A_2 & \vdots \\ \vdots & \ddots & \circledast \\ \circledast & \dots & A_\ell \end{pmatrix} = A_1 \oplus A_2 \oplus \dots \oplus A_\ell$$

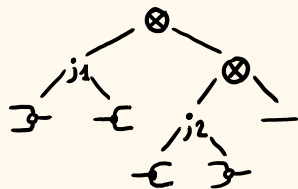
THEOREM

$$\max_i \text{rank } A_i \leq \text{mwd } A \leq \max_i \text{rank } A_i + 1$$

# MONOIDAL WIDTH OF MATRICES - EXAMPLE

$$\mathcal{M} = \{ \text{---}, \text{X}, \text{---} \circ \text{---}, \text{---} \circ \text{---}, \text{---} \text{---} \text{---}, \text{---} \bullet \}$$

$$A = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \text{---} \circ \text{---} \text{---} \text{---}$$



THEOREM

$$\max_i \text{rank } A_i \leq \text{mwd } A \leq \max_i \text{rank } A_i + 1$$

# OUTLINE

- monoidal decompositions
- monoidal width for matrices
- monoidal width for graphs

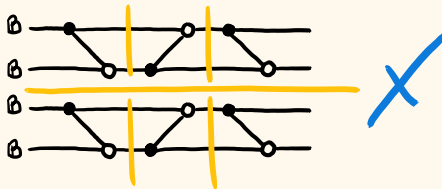
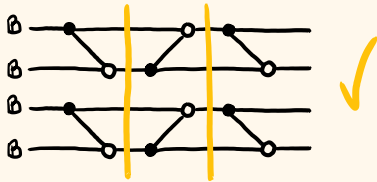


# VARIATIONS OVER MONOIDAL WIDTH (1)

$(\mathcal{C}, \mathcal{F}, \mathcal{O}, w)$  decomposition system  
generators of  $\mathcal{C}$   $\rightarrow$  weight function  
allowed operations

## CATEGORICAL PATH WIDTH

$$\mathcal{O} = \{ ;_x \text{ for any object } x \}$$



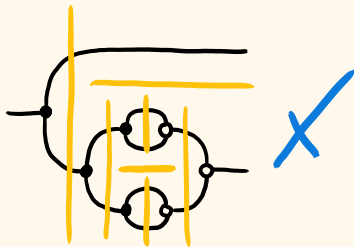
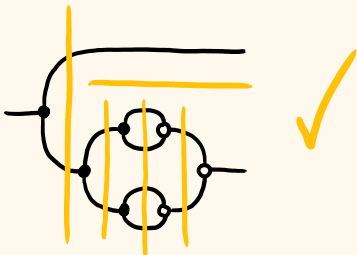
# VARIATIONS OVER MONOIDAL WIDTH (2)

$(\mathcal{C}, \mathcal{F}, \mathcal{O}, w)$  decomposition system  
generators of  $\mathcal{C}$   $\rightarrow$  weight function  
allowed operations

## CATEGORICAL TREE WIDTH

$\mathcal{O} = \{ \otimes, ;_x \text{ for any object } x \}$

composition restricted to generators on the left



# TREE WIDTH [Robertson & Seymour, 1986]

$G = (V, E)$  undirected graph

TREE DECOMPOSITION

$(T, \lambda)$  where

- $T$  binary tree
- $\lambda : \text{nodes}(T) \rightarrow \text{subgraphs}(G)$  such that

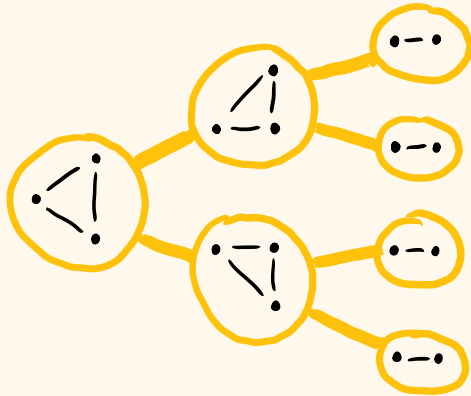
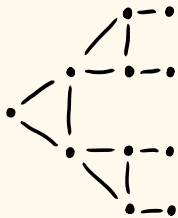
$$\left\{ \begin{array}{l} \bigcup_{m \in \text{nodes}(T)} \lambda(m) = G \\ \forall m \in \text{nodes}(T) \quad \lambda(m) \supseteq \lambda(m.\text{left}) \cup \lambda(m.\text{right}) \end{array} \right.$$

$$\forall m \in \text{nodes}(T) \quad \lambda(m) \supseteq \lambda(m.\text{left}) \cup \lambda(m.\text{right})$$

# TREE WIDTH - EXAMPLE

$G = (V, E)$  undirected graph

TREE DECOMPOSITION



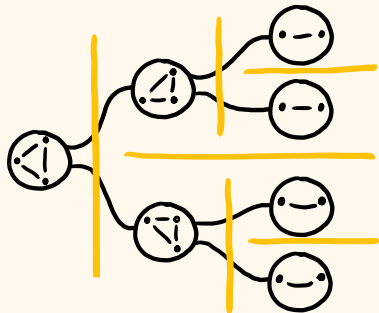
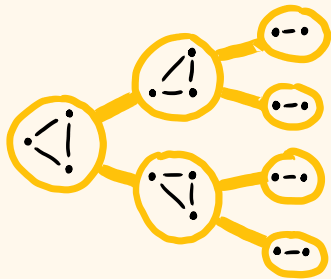
# TREE WIDTH & MONOIDAL WIDTH

$G = (V, E)$  undirected graph

$\mathcal{G} = \emptyset \rightarrow G \leftarrow \emptyset$  cospan corresponding to  $G$

THEOREM

$$\text{twid } G \leq \text{mwid}_T \mathcal{G} \leq 2 \cdot \text{twid } G$$



# SUMMARY OF RESULTS

- matrices  $\text{mwd} \Leftrightarrow \text{rank}$
- cospanns of graphs
  - $\text{mwd}_T \Leftrightarrow \text{tree width}$
  - $\text{mwd}_p \Leftrightarrow \text{path width}$
  - $\text{mwd} \Leftrightarrow \text{branch width}$
- prop of graphs  $\text{mwd} \Leftrightarrow \text{rank width}$

# FUTURE WORK

- monoidal width in other categories
- directed tree width, DAG width, Kelly width, cut width, clique width, ...
- monoidal width 'functorially'  
(describe the weight function as a functor)

THANKS!



# SOME REFERENCES

## GRAPH WIDTHS

- Robertson & Seymour, Graph minors 1-X, 1983-1991
- Oum & Seymour, Approximating clique width and branch width, 2006
- Johnson, Robertson, Seymour & Thomas, Directed tree width, 2001
- Courcelle & Oularin, Upper bounds to the clique width of graphs, 2000
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